

70: Discrete Math and Probability Theory

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Questions/Announcements

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Questions/Announcements \implies piazza:

piazza.com/berkeley/fall2017/cs70

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Programming + Microprocessors

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What can computers do?

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Probability!

See note 1, for more discussion.

Wason's experiment:1

Suppose we have four cards on a table:

- ▶ 1st about Alice, 2nd about Bob, 3rd Charlie, 4th Donna.

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Suppose we have four cards on a table:

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- ▶ Consider the theory:
"If a person travels to Chicago, he/she flies."
- ▶ Suppose you see that Alice went to Baltimore, Bob drove, Charlie went to Chicago, and Donna flew.

Alice
Baltimore

Bob
drove

Charlie
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- ▶ Which cards must you flip to test the theory?

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Answer:

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- ▶ Which cards must you flip to test the theory?

Answer: Later.

CS70: Lecture 1. Outline.

Today: Note 1.

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The language of proofs!

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Today: Note 1. Note 0 is background. Do read it.

The language of proofs!

1. Propositions.
2. Propositional Forms.
3. Implication.
4. Truth Tables
5. Quantifiers
6. More De Morgan's Laws

Propositions: Statements that are true or false.

$\sqrt{2}$ is irrational

$$2+2 = 4$$

$$2+2 = 3$$

826th digit of pi is 4

Johny Depp is a good actor

All evens > 2 are sums of 2 primes

$$4 + 5$$

$$x + x$$

Alice travelled to Chicago

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Propositional Forms.

Put propositions together to make another...

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Conjunction (“and”): $P \wedge Q$

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$\neg (2 + 2 = 4)$ – a proposition that is ...

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“ $2 + 2 = 3$ ” \wedge “ $2 + 2 = 4$ ” – a proposition that is ...

Propositional Forms.

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\neg “ $(2 + 2 = 4)$ ” – a proposition that is ... **False**

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Propositional Forms: quick check!

$P = \text{"}\sqrt{2} \text{ is rational"}$

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$P = \text{"}\sqrt{2} \text{ is rational"}$

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$P \wedge Q \dots$

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$P \vee Q$... True

$\neg P$... True

Put them together..

Propositions:

P_1 - Person 1 rides the bus.

Put them together..

Propositions:

P_1 - Person 1 rides the bus.

P_2 - Person 2 rides the bus.

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Put them together..

Propositions:

P_1 - Person 1 rides the bus.

P_2 - Person 2 rides the bus.

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But we can't have either of the following happen; That either person 1 or person 2 ride the bus and person 3 or 4 ride the bus. Or that person 2 or person 3 ride the bus and that either person 4 rides the bus or person 5 doesn't.

Put them together..

Propositions:

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$\neg(((P_1 \vee P_2) \wedge (P_3 \vee P_4)) \vee ((P_2 \vee P_3) \wedge (P_4 \vee \neg P_5)))$

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Can person 3 ride the bus?

Put them together..

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Can person 3 ride the bus?

Can person 3 and person 4 ride the bus together?

Put them together..

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Can person 3 ride the bus?

Can person 3 and person 4 ride the bus together?

Put them together..

Propositions:

P_1 - Person 1 rides the bus.

P_2 - Person 2 rides the bus.

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But we can't have either of the following happen; That either person 1 or person 2 ride the bus and person 3 or 4 ride the bus. Or that person 2 or person 3 ride the bus and that either person 4 rides the bus or person 5 doesn't.

Propositional Form:

$\neg(((P_1 \vee P_2) \wedge (P_3 \vee P_4)) \vee ((P_2 \vee P_3) \wedge (P_4 \vee \neg P_5)))$

Can person 3 ride the bus?

Can person 3 and person 4 ride the bus together?

This seems ...

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We need a way to keep track!

Truth Tables for Propositional Forms.

P	Q	$P \wedge Q$
T	T	T
T	F	
F	T	
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Notice: \wedge and \vee are commutative.

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One use for truth tables: Logical Equivalence of propositional forms!

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T	F	F	F
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DeMorgan's Law's for Negation: distribute and flip!

$$\neg(P \wedge Q)$$

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F	F	T	T

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$$\neg(P \wedge Q) \equiv \neg P \vee \neg Q$$

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$$\neg(P \wedge Q) \quad \equiv \quad \neg P \vee \neg Q \qquad \neg(P \vee Q)$$

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T	T	F	F
T	F	F	F
F	T	F	F
F	F	T	T

DeMorgan's Law's for Negation: distribute and flip!

$$\neg(P \wedge Q) \equiv \neg P \vee \neg Q \qquad \neg(P \vee Q) \equiv \neg P \wedge \neg Q$$

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

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$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q,$

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$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R)$$

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

$$\text{RHS: } (T \wedge Q) \vee (T \wedge R)$$

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

$$\text{RHS: } (T \wedge Q) \vee (T \wedge R) \equiv (Q \vee R).$$

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

$$\text{RHS: } (T \wedge Q) \vee (T \wedge R) \equiv (Q \vee R).$$

P is False .

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

$$\text{RHS: } (T \wedge Q) \vee (T \wedge R) \equiv (Q \vee R).$$

P is False .

$$\text{LHS: } F \wedge (Q \vee R)$$

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

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$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

$$\text{RHS: } (T \wedge Q) \vee (T \wedge R) \equiv (Q \vee R).$$

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Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

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Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

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P is False .

$$\text{LHS: } F \wedge (Q \vee R) \equiv F.$$

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Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

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Distributive?

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Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

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Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

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P is True .

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$$P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)?$$

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

$$\text{RHS: } (T \wedge Q) \vee (T \wedge R) \equiv (Q \vee R).$$

P is False .

$$\text{LHS: } F \wedge (Q \vee R) \equiv F.$$

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Simplify: $T \vee Q \equiv T$,

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

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P is False .

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$$P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)?$$

Simplify: $T \vee Q \equiv T$, $F \vee Q \equiv Q$.

Distributive?

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$$P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)?$$

Simplify: $T \vee Q \equiv T$, $F \vee Q \equiv Q$.

Foil 1:

Distributive?

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Simplify: $T \vee Q \equiv T$, $F \vee Q \equiv Q$.

Foil 1:

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Simplify: $T \vee Q \equiv T$, $F \vee Q \equiv Q$.

Foil 1:

$$(A \vee B) \wedge (C \vee D) \equiv (A \wedge C) \vee (A \wedge D) \vee (B \wedge C) \vee (B \wedge D)?$$

Foil 2:

Distributive?

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)?$$

Simplify: $(T \wedge Q) \equiv Q$, $(F \wedge Q) \equiv F$.

Cases:

P is True .

$$\text{LHS: } T \wedge (Q \vee R) \equiv (Q \vee R).$$

$$\text{RHS: } (T \wedge Q) \vee (T \wedge R) \equiv (Q \vee R).$$

P is False .

$$\text{LHS: } F \wedge (Q \vee R) \equiv F.$$

$$\text{RHS: } (F \wedge Q) \vee (F \wedge R) \equiv (F \vee F) \equiv F.$$

$$P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)?$$

Simplify: $T \vee Q \equiv T$, $F \vee Q \equiv Q$.

Foil 1:

$$(A \vee B) \wedge (C \vee D) \equiv (A \wedge C) \vee (A \wedge D) \vee (B \wedge C) \vee (B \wedge D)?$$

Foil 2:

$$(A \wedge B) \vee (C \wedge D) \equiv (A \vee C) \wedge (A \vee D) \wedge (B \vee C) \wedge (B \vee D)?$$

Implication.

$P \implies Q$ interpreted as

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If P , then Q .

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True Statements: $P, P \implies Q$.

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Examples:

Statement: If you stand in the rain, then you'll get wet.

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Conclude: Q is true.

Examples:

Statement: If you stand in the rain, then you'll get wet.

P = "you stand in the rain"

Implication.

$P \implies Q$ interpreted as

If P , then Q .

True Statements: $P, P \implies Q$.

Conclude: Q is true.

Examples:

Statement: If you stand in the rain, then you'll get wet.

P = "you stand in the rain"

Q = "you will get wet"

Implication.

$P \implies Q$ interpreted as

If P , then Q .

True Statements: $P, P \implies Q$.

Conclude: Q is true.

Examples:

Statement: If you stand in the rain, then you'll get wet.

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Statement: "Stand in the rain"

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Statement: If you stand in the rain, then you'll get wet.

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Statement:

If right triangle has sidelengths $a, b \leq c$, then $a^2 + b^2 = c^2$.

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Statement: If you stand in the rain, then you'll get wet.

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Statement: If you stand in the rain, then you'll get wet.

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If right triangle has sidelengths $a, b \leq c$, then $a^2 + b^2 = c^2$.

P = "a right triangle has sidelengths $a \leq b \leq c$ ",

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Non-Consequences/consequences of Implication

The statement " $P \implies Q$ "

Non-Consequences/consequences of Implication

The statement " $P \implies Q$ "

only is **False** if P is **True** and Q is **False** .

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Some Fun: use propositional formulas to describe implication?

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Some Fun: use propositional formulas to describe implication?

$((P \implies Q) \wedge P) \implies Q$.

Implication and English.

$$P \implies Q$$

- ▶ If P , then Q .

Implication and English.

$$P \implies Q$$

- ▶ If P , then Q .
- ▶ Q if P .
Just reversing the order.

Implication and English.

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Implication and English.

$$P \implies Q$$

- ▶ If P , then Q .
- ▶ Q if P .
Just reversing the order.
- ▶ P only if Q .
Remember if P is true then Q must be true.
this suggests that P can only be true if Q is true.
since if Q is false P must have been false.

Implication and English.

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This means that proving P allows you
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since if Q is false P must have been false.
- ▶ P is sufficient for Q .
This means that proving P allows you
to conclude that Q is true.
- ▶ Q is necessary for P .
For P to be true it is necessary that Q is true.
Or if Q is false then we know that P is false.

Truth Table: implication.

P	Q	$P \implies Q$
T	T	T
T	F	
F	T	
F	F	

Truth Table: implication.

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T	T	T
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Truth Table: implication.

P	Q	$P \implies Q$
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T	F	F
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P	Q	$\neg P \vee Q$
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T	F	
F	T	
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Truth Table: implication.

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$$\neg P \vee Q \equiv P \implies Q.$$

These two propositional forms are logically equivalent!

Contrapositive, Converse

- ▶ Contrapositive of $P \implies Q$ is $\neg Q \implies \neg P$.

Contrapositive, Converse

- ▶ Contrapositive of $P \implies Q$ is $\neg Q \implies \neg P$.
 - ▶ If the plant pollutes, fish die.

Contrapositive, Converse

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(contrapositive)

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$$P \implies Q$$

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If fish die the plant pollutes.

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If fish die the plant pollutes.
Not logically equivalent!

Contrapositive, Converse

- ▶ Contrapositive of $P \implies Q$ is $\neg Q \implies \neg P$.
 - ▶ If the plant pollutes, fish die.
 - ▶ If the fish don't die, the plant does not pollute.
(contrapositive)
 - ▶ If you stand in the rain, you get wet.
 - ▶ If you did not stand in the rain, you did not get wet.
(not contrapositive!) converse!
 - ▶ If you did not get wet, you did not stand in the rain.
(contrapositive.)

Logically equivalent! Notation: \equiv .

$$P \implies Q \equiv \neg P \vee Q \equiv \neg(\neg Q) \vee \neg P \equiv \neg Q \implies \neg P.$$

- ▶ Converse of $P \implies Q$ is $Q \implies P$.
If fish die the plant pollutes.
Not logically equivalent!
- ▶ **Definition:** If $P \implies Q$ and $Q \implies P$ is P if and only if Q or $P \iff Q$.
(Logically Equivalent: \iff .)

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Next: Statements about boolean valued functions!!

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Wait! What is \mathbb{N} ?

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- ▶ See note 0 for more!

Lecture ended here. Will continue from here on Friday.